

Understanding, Predicting and Scheduling Serverless Workloads under *Partial Interference*

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Serverless Computing



Serverless Workloads

Scheduled Background (BG)

IoT data collection, monitoring etc.

Short-term Computing (SC)

Mapreduce/Spark, linear algebra etc:

Latency Sensitive (LS)

Search, e-commerce, social network.





 ≤ 90

Serverless Workloads



• A server can accommodate hundreds of functions.



Functions are started or released at all

times.

interference

Interference



Interference causes high tail latency.

Cutting Tail Latency

Physical isolation



Software optimization

VM + Container



Low resource efficiency

Overprovisioning

Cutting Tail Latency



 Only fine-grained and proactive control can provide good performance and high throughput for serverless.

Partial Interference



 Interference occurs only at some parts, but not all, of the workloads.

Partial Interference in Serverless



Partial Interference in Serverless

• LS: Social network (9 functions)



- BG/SC:
 - matrix multiplication,
 - video processing,
 - iperf,
 - dd,
 - LR,
 - Kmeans etc.

Serverless makes partial interference particularly severe.

High volatility

 Serverless functions are diverse in terms of execution behavior and resource consumption, making partial interference more volatile.

Partial interference scenarios



- Either as strong as *Full interference*, or as weak as Zero interference.
- The difference in the 99th percentile latency among these scenarios reaches 7×.

Spatial variation

 Serverless functions are small in size and stateless, making partial interference spatially varied.

Partial interference scenarios



• Interference on the critical path (e.g., 1-2-6-8-9) generates a much more severe impact than interference on the non-critical path.

3

Temporal variation

 Serverless functions are short-lived, making partial interference temporally varied.



• The maximum difference in JCT of colocated Logistic Regression (LR) and Kmeans is more than 2×.

4 Hotspot propagation

 Partial interference triggers a chain reaction across the function call path and leads to diametrically opposite effects.



- The QPS of the subsequent invoked functions decreases.
- The bottlenecked gateway degrades the invocation speeds of all other functions.

5 Restoring propagation

 The local control of partial interference suffers from impact propagation.



(b) Inter. at <a>(c) compose-and-upload

• Local interference control increase the other functions' latencies due to the restored invocations.

Predicting

6 Predictability in serverless

 Accurate partial interference prediction is enabled by function-level profiles, thereby improving the QoS of workloads.



Function-level profiles

produce an average median that is 2× lower than that by workload-level profiles.

Predicting



DA "spatial-temporal interference" - aware incremental learning predictor, which can converge quickly by training on the profiles of functions along an end-to-end call path.







Function-level profiling

- solo-run way
- non-intrusive
 - system-layer
 - microarchitecture-layer

In the second second



 $P_{A\cup\{B,C,...\}} = RM(R_A, R_B, R_C, ..., U_A, U_B, U_C, ..., D_A, D_B, D_C, ..., T_A, T_B, T_C, ...)$

AllocatedRes Utilization

Lifetime

of rows = # of servers



- $(D_2, D_3) = (0, t_{delay})$
- U₁, U₂.

 u_{ij}^k : the *k*th metric measured when workload *i* is deployed on server *j*.

Learning models

- IKNN, IRFR, IMLP, ILR, ISVR, ESP [Mishra2017], Pythia [Xu2018]
- The prediction error of IPC generated by **IRFR** is as low as **1.71%**.



Figure 9. The prediction errors of (a) IPC and (b) tail latency.

Scheduling

- A simple binary search scheduling algorithm
 - maximize resource efficiency, by deploying function instances on a min number of active servers, while guaranteeing the QoS of colocated workloads.



- Workloads: [Gsight and data available: https://github.com/tjulym/gray]
 - BG/SC: ServerlessBench [Yu2020], FunctionBench [Kim2019];
 - LS: social network, e-commerce;
 - Azure trace [Shahrad2020]
- Testbed:

Openfaac	Component	Specification	Component	Specification
Openiaas	CPU model	Intel Xeon E7-4820v4	Shared LLC Size	25MB
	Number of sockets	4	Memory Capacity	256GB
	Processor BaseFreq.	2.00 GHz	Operating System	Ubuntu 14.04.5LTS
	Threads	80 (40 physical cores)	SSD Capacity	960GB
	Private L1&L2 Cache	64 KB and 256 KB	Number of Nodes	8

• Fast and stable convergence:

- Function-level profiles enable Gsight to converge quickly. Its precision is also stable after convergence.
- To achieve the same prediction error, Gsight-IRFR only requires 1/3 samples compared to serverful system.



Function density



Gsight scheduling can improve the function density by an average of 18.79% and by 48.48% over those of *Pythia* and *Worst Fit*.

CPU utilization



Gsight scheduling can improve CPU utilization by **30.02%** and **67.51%** on average compared with that of *Pythia* and *Worst Fit*, respectively.

SLA guarantee



 Gsight scheduling can guarantee the SLA of social network 95.39% of the time.

Overhead



• Training:

- 3,000 samples takes us less than 2 person-hours, the training process takes only < 10 minutes.
- Gsight predicting takes 3.48 ms.
 - Cold start is slow.
 - A scalable gateway is also required.

Conclusion

- From interference to partial interference, we are moving forward to understand more about tail latency.
- Serverless makes the resource management more challengeable.
- Only fine-grained and proactive control can provide good performance and high throughput for serverless.
- Gsight is just a start!

